Subject: Re: Kqueue/epoll based interface for TcpSocket and WebSocket Posted by mirek on Mon, 30 Apr 2018 09:20:50 GMT

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```
Sorry for the delay...
Quote:
template <class T SOCKET>
class SocketEventQueue: NoCopy
{
public:
SocketEventQueue(): errorCode(NOERR) { InitEventQueue(); }
~SocketEventQueue():
bool ClearEventQueue();
QueueHandler GetQueueHandler() const:
bool IsError() const { return errorCode != NOERR; }
ErrorCode GetErrorCode();
bool SubscribeSocketRead(const T SOCKET &sock);
bool SubscribeSocketWrite(const T SOCKET &sock):
bool SubscribeSocketReadWrite(const T SOCKET &sock):
bool DisableSocketRead(const T SOCKET &sock):
bool DisableSocketWrite(const T_SOCKET &sock);
bool DisableSocketReadWrite(const T SOCKET &sock);
bool RemoveSocket(const T_SOCKET &sock);
bool IsSocketSubscribedRead(const T_SOCKET &sock) const { return IsSockeSubscribed(sock,
WAIT READ); }
bool IsSocketSubscribedWrite(const T_SOCKET &sock) const { return IsSockeSubscribed(sock,
WAIT_WRITE); }
bool IsSocketDisabledRead(const T_SOCKET &sock) const { return IsSockeDisabled(sock,
WAIT READ); }
bool IsSocketDisabledWrite(const T_SOCKET &sock) const { return IsSockeDisabled(sock,
WAIT WRITE); }
Vector<SocketEvent<T_SOCKET>> Wait(int timeout);
};
```

What is DisableSocketRead supposed to do? Opposite of Subscribe?

If I am right about T_SOCKET, I have to disagree with the interface a bit.

Particulary, I think

Vector<SocketEvent<T_SOCKET>> Wait(int timeout);

is clumsy - this will IMO cause problems with mapping T_SOCKET back to its "processes".

When I was thinking about how to proceed with this, I was considering to simply expand SocketWaitEvent. I believe that the best would probably be to treat it as full array of sockets, using indices to identify the 'process'. Something like

```
class SocketWaitEvent {
public:
void Clear()
                                         { socket.Clear(); }
void Add(SOCKET s, dword events)
                                                    { socket.Add(MakeTuple((int)s, events)); }
void Add(TcpSocket& s, dword events)
                                                     { Add(s.GetSOCKET(), events); }
int Wait(int timeout);
dword Get(int i) const:
dword operator[](int i) const
                                             { return Get(i); }
// new:
     void Set(int ii, TcpSocket& s, dword events);
     void Insert(int ii, TcpSocket& s, dword events);
     void Remove(int ii, TcpSocket& s, dowrd events);
     Vector<int> WaitEvent(int timeout);
 // or perhaps
     Vector<Tuple<int, dword>> WaitEvent(int timeout); // dword part contains Get bitmask
// maybe:
     void Clear(int ii); // makes index empty
     int FindEmpty() const; // finds the first index that is empty
SocketWaitEvent();
};
```

Is there a reason to make things more complicated that this?

Quote:

The first problem is related to the current implementation of TcpSocket::RawWait(...). It uses

select(...) system call for determining possibility of reading/writing data or exceptional state of socket. As far as I can understand, it means that server with large nubmer of sockets will work slowly anyway. I patched TcpSocket::RawWait(...) for BSD platform on my local machine (see below) before I started to implement the interface. Now I think that it's possible to use SocketEventQueue in purpose of determining socket state instead of raw kqueue/epoll/select. What do you think about this?

I agree.

Quote:

But there is another problem related to kqueue/epoll reaction on socket closing for both my patch and SocketEventQueue. So here's the patch:

```
#ifdef PLATFORM BSD
timespec *tvalp = NULL;
timespec tval;
if(end_time != INT_MAX || WhenWait) {
 to = max(to, 0);
 tval.tv sec = to / 1000;
 tval.tv_nsec = 1000000 * (to % 1000);
 tvalp = &tval;
 if (to)
 LLOG("RawWait timeout: " << to);
struct kevent eventrx, eventw;
struct kevent triggeredEvents[2];
int ka:
int eventFlags = EV_ADD | EV_ONESHOT;
if( (kq = kqueue()) == -1) // queue fd should be created once at the moment of socket opening
      // and closed at the moment of socket closing
     // the same is for SocketEventQueue object
 LLOG("kg = kgueue() returned -1");
 SetSockError("wait");
 return false:
}
if(flags & WAIT READ)
 EV_SET( &eventrx, socket, EVFILT_READ, eventFlags, 0, 0, NULL );
 if( kevent( kg, &eventrx, 1, NULL, 0, NULL ) == -1 )
 LLOG("kevent( kg, &eventrx, 1, NULL, 0, NULL ) returned -1");
```

```
SetSockError("wait");
 close(kg);
 return false;
if(flags & WAIT_WRITE)
 EV_SET( &eventw, socket, EVFILT_WRITE, eventFlags, 0, 0, NULL );
 if( kevent( kg, &eventw, 1, NULL, 0, NULL ) == -1 )
 LLOG("kevent( kg, &eventw, 1, NULL, 0, NULL ) returned -1");
 SetSockError("wait");
 close(kq);
 return false;
 }
}
int avail = kevent( kq, nullptr, 0, triggeredEvents, 2, tvalp ); // here is the problem if
socket
                                              // works in blocking mode
                                              // or if timeout is too long
close(kq);
#else
  // default select implementation
Now let's imagine the situation:
TcpSocket server; // passes through TcpSocket::Listen()
void Server() // runs in several threads
static StaticMutex serverMutex;
while(!Thread::IsShutdownThreads())
 TcpSocket
               client;
 bool acceptStatus;
 Mutex::Lock __(serverMutex);
 //acception is in blocking mode
 acceptStatus = client.Accept(server); // calls TcpSocket::RawWait(...)
 }
```

```
... // connection handling
}

...

void SignalHandler(int sig)
{
    server.Close(); // Doesn't interrupt kevent system call in TcpSocket::RawWait(...)
    // close(socket) just makes kqueue to delete all events
    // associated with socket descriptor from it's kernel queue

Thread::ShutdownThreads();
}
```

So I can't normally terminate the server if I work with sockets in blocking mode.

Do you have any ideas how to interrupt kevent waiting loop?

I've tried to call shutdown(socket, SD_BOTH) for sockets that hadn't been passed through TcpSocket::Listen(), and it works for me.

But I still can't deal with listening socket. Solution I've found is to use pipe-trick: read-end descriptor attaches to kqueue/epoll, and write-end descriptor attaches to socket. When socket closes, it writes some data in pipe with write-end descriptor.

But it means that socket must hold all queue write-end descriptors it was attached.

Could you help me with my problem?[/quote]

Well, I was fighting with this one too, years ago. Thats nasty little problem there.

I the end, I believe that the best solution is to make Accept return until there are any active threads by doing localhost connect.

```
TcpSocket s;
s.Connect("127.0.0.1", port);
```

You can check Skylart/App.cpp.

Another option, not always applicable, is not to bother and let the signal kill the application just as it is supposed to....:)